Performance evaluation of RYU controller under distributed denial of service attacks

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Article Info

Article history:

Received Jul 6, 2022 Revised Jun 27, 2023 Accepted Jul 2, 2023

Keywords:

CPU consumption Distributed denial of service Link latency RAM exhaustion RYU controller Software-defined networking

ABSTRACT

Distributed denial of service (DDoS) attacks have been identified as one of the greatest threats to software-defined networking (SDN) because they are highly effective, hard to detect, and easy to use, and they take advantage of vulnerabilities in which the new architecture still exists. This paper describes one technique for denying the RYU controller's services, which can cause the controller's resources to be depleted if a significant number of packets from various zombie hosts are sent to the controller using spoofed source internet protocol (IP) addresses. In order to demonstrate the impact of the attack, we measure various metrics related to RYU controllers such as its central processing unit (CPU) usage and network throughput. In this work, Mininet was used to simulate the data plane and measure metrics such as random access memory (RAM) usage, CPU load, and link latency.

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1. INTRODUCTION

Software-defined networking has drawn a lot of interest as a way to get around the drawbacks of conventional distributed networks. It has a number of benefits, but the separation of the control plane and data plane-which makes the network infrastructure more flexible and controllable-is its main advantage [1], [2]. Although software-defined networking (SDN) represents a promising future network architecture, distributed denial of service (DDoS) attacks can exploit its centralized configuration settings. Computer networks are significantly at risk from DDoS attacks because they are difficult to detect due to their distributed nature and relatively easy to carry out [3]. Because of this flaw, incoming packets from switches will flood the SDN controller, causing it to overload [4]. The three layers that make up an SDN are the data plane (data plane), the control plane (control plane), and the application plane (apps plane) (management plane) as shown in Figure 1. The control plane is a controller that oversees the network as a whole. Decisions about the management, quality of service (QoS), and security of the network are all made by the controller [5]–[7].

Using open, standardized interfaces known as south-bound interfaces and protocols like OpenFlow, the SDN controller connects with switches (forwarding only devices) in the data plane. To carry out a variety of tasks, including routing, switching, firewalling, network address translation, and load balancing, the controller directs the forwarding elements (data plane switches) [8]. SDN offers an innovative approach to network management. In SDN, switches do not process incoming packets; instead, they merely search their forwarding tables for a match. If no match is found, the packets are forwarded to the controller for processing.

SDN's controller is its operating system. It processes the packets and determines whether they will be forwarded or discarded by the switch. Using this method, SDN divides the forwarding and processing planes [7].

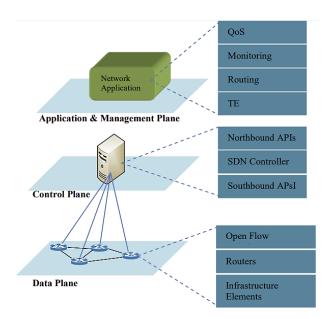


Figure 1. SDN architecture [6]

While this new architecture represents significant gains [9], it is not without flaws and security risks [10]. Several problems have been highlighted, including an absence of authentication measures for packet senders [11]. Due to these flaws, a number of methods have been used to target the controller and endpoints' accessibility, integrity, and confidentiality. DoS and DDoS attacks, in particular the disruption of the controller to prevent switches from processing requests, have been modified to suit the SDN architecture.

A significant risk is posed by the centralized nature of network control in SDN. The effects of these attacks on the controller, if it is compromised, will be felt throughout the entire network. Compared to traditional networks, where control is distributed, this amplifies the effects of such attacks. Consequently, a controller attack in an SDN environment can have a much bigger impact than a comparable attack on a traditional environment.

DoS and DDoS [12], [13] attacks have reportedly been undertaken against a number of SDN controllers [14], including FloodLight, NOX, and RYU [15]-[17], among others, according to past studies. To the best of our knowledge, no in-depth research has been conducted on DoS and DDoS attacks conducted via POX and RYU controllers. Therefore, the primary objective of this research is to look at the impact of a DDoS attacks on the RYU controller.

2. RELATED WORK

Alrammahi and Bhaya [18] employed a Mininet emulator and an OpenFlow switch connected to a POX controller to conduct an experimental test. The aim was to implement and compare the performance of four load balancing algorithms: random, round robin (RR), weighted round robin (WRR), and least connections (LC). Various parameters were considered to evaluate the effectiveness of the comparison process. These included workload, load balancing level, and time-related values such as max-time, min-time, average-time, and response time. These parameters were analyzed to assess the performance and efficiency of the load balancing algorithms being tested.

Mahdi and Abdullah [19] suggest network infrastructure configured using a Mininet network emulator, and the FlowVisor hypervisor is used to virtualize the infrastructure and enable multiple SDN controllers. Additionally, three scenarios are presented for network slicing implementation and comparison. The findings indicate that as the number of tenants increases and the network is divided into more than two slices, the performance of the network approaches and may match that of the network without the use of FlowVisor, which results in the control of the delay rate and the data transfer rate.

Noman and Jasim [20] uses network monitoring tools like Iperf and distributed internet traffic generator (D-ITG) to measure and record performance indicators including service delay, utilized bandwidth,

received packets and bytes in order to examine the controller's capabilities and assess how well they function in an SDN setting. Bhardwaj and Panda [21] describes an implementation of SDN architecture for network traffic analysis using an open-source RYU SDN controller. The proposed work examines the performance of a custom network topology based on SDN architecture for node-to-node performance characteristics such as bandwidth, throughput, and round-trip time.

Abood and Abdul-Majeed [22] employs the entropy-threshold technique, which is programmed in Python within the controllers, as well as DDoS attacks carried out by internet of things (IoT) sensors. If the entropy value falls below a certain threshold, the controller will block that specific port in the switch and shut it down. The results show that FloodLight outperforms RYU controllers in dealing with DDoS attack cases. In addition, the FloodLight controller's Throughput decreased from 0.1158 KB/s to 0.08916 KB/s, while the RYU controller's throughput decreased from 0.0646 KB/s to 0.03294 KB/s.

3. SDN COMPONENTS

SDN can be identified by its five main attributes: separation of planes, streamlined devices, centralized control, network automation, and virtualization, all of which promote openness [22]. OpenFlow is essential component in SDN architecture. It is uses to packet forwarding or flow statistics collection. OpenFlow switches, a secure channel, and a controller [23] make up the bulk of the architecture. Flow tables are maintained by the devices, better known as switches. Flow, counters, and actions can be uniquely identified for each entry using the several match fields that are provided for this purpose. It is possible to define a flow as a group of packets that are all in the same time period and have the same properties. An optional encryption technology known as transport layer security (TLS) can be used to encrypt data between data plane and control plane [24]. The controller is in charge of maintaining the flow tables on the switch. It is possible for the controller to make OpenFlow requests or add/remove entries from the flow tables.

Figure 2 depicts the packet forwarding mechanism in an OpenFlow switch. The header packet are captured and inspect them when the packet arrives at the switch for matching reasons. To find a match, the flow rules recorded in the flows table are compared to the header packet in ascending order using the flow rules. To begin, the matching system looks for a flow rule that matches this packet to an existing defined flow; if it does, the switch performs the action described in this rule. If a match is not found, the switch delivers the packet information to the controller in an OpenFlow packet called packet in, which is then processed. Once this flow has been identified, the controller creates a new switch rule for this flow [23].

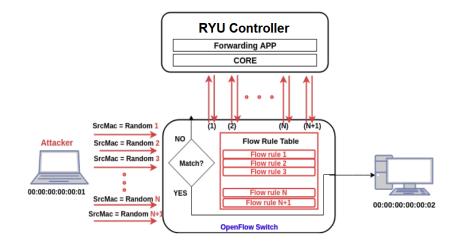


Figure 2. Packets forwarding in OpenFlow switch

4. LAUNCH DDOS ATTACK IN SDN

The RYU controller is the target of a DDoS attack, which is DoS initiated by many compromised hosts (zombies) known as zombies [25]. To carry out a DoS attack, the attacker takes advantage of the controller's limited resources. The DDoS strategy entails flooding the controller with a large number of fabricated packets in order to deplete its available resources, such as central processing unit (CPU) processing power and network bandwidth as shown in Figure 2. During our testing, we discovered that as each packet is processed, the RYU forwarding function generates a new flow rule with the source and destination MAC

addresses as matching criteria. The DoS attack is based on a heuristic that focuses on overwhelming the controller by sending a large number of fabricated packets in a continuous stream. This places an undue strain on the controller's resources, as each packet causes the RYU forwarding application to create a new flow rule. This accumulation of flow rules and associated processing tasks can eventually exhaust the controller's CPU and network bandwidth, resulting in severe performance degradation and a negative impact on overall network operation.

It is common knowledge that a single Python script was written utilizing the Scapy package. It's an advanced user interface for manipulating packets. It can send, receive, alter, and decode packets of many different protocols, as well as match requests with their responses. Scanning, tracing, probing, unit testing, assaults, and network discovery are just some of the commonplace activities it can accomplish with ease (it can replace hping, 85% of nmap, arpspoof, arp-sk, arping, tcpdump, tshark, and p0f,) [26] as shown in Figure 3.

When the fabricated packets reach any switch in this attack scenario, they must be forwarded to the controller as a single packet. The attacker sends packets from multiple compromised hosts, with the goal of exceeding the controller's threshold by flooding it with a high influx of fabricated packets every second. In this attack, internet control message protocol (ICMP) packets are preferred over user datagram protocol (UDP) packets. The reason for this selection is that the Scapy tool used for packet creation and transmission spends less time generating and transmitting ICMP packets, allowing for a higher transmission rate.

X "Node: h1"@mininet-vm <Ether type=0x800 I<IP frag=0 proto=udp src=29.120.184.5 dst=['10.0.0.8'] I<UD</pre> sport=http dport=1 I>>> Sent 1 packets. 139.54.218.168 KEther type=0x800 |<IP frag=0 proto=udp src=139.54.218.168 dst=['10.0.0.8'] |<UIP sport=http dport=1 |>>> Sent 1 packets. 32.44.180.176 <Ether type=0x800 |<IP frag=0 proto=udp src=32,44,180,176 dst=['10,0,0,8'] |<UDP sport=http dport=1 |>>> Sent 1 packets. 25.189.160.164 <Ether type=0x800 I<IP frag=0 proto=udp src=25.189.160.164 dst=['10.0.0.8'] I<UIP sport=http dport=1 I>>> Sent 1 packets. 14.158.123.252 <Ether type=0x800 |<IP frag=0 proto=udp src=14.158.123.252 dst=['10.0.0.8'] |<UDP sport=http dport=1 |>>> Sent 1 packets. 144.77.246.199 <Ether type=0x800 I<IP frag=0 proto=udp src=144.77.246.199 dst=['10.0.0.8'] I<UDP sport=http dport=1 I>>> . Sent 1 packets. 62.225.40.89 <Ether type=0x800 |<IP frag=0 proto=udp src=62.225.40.89 dst=['10.0.0.8'] |<UDP sport=http dport=1 |>>> Sent 1 packets. 210.106.82.137 <Ether type=0x800 I<IP frag=0 proto=udp src=210.106.82.137 dst=['10.0.0.8'] I<UIP sport=http dport=1 I>>>

Figure 3. Traffic generation using Scapy

5. EVALUATION

It will be clarified as SDN environment and the remaining components where it is a type of flood attack, such as the Scapy tool, to display the outcomes of the testing phase. Multiple packets are sent to the network device in order to interrupt or slow down the service. The CPU and random-access memory (RAM) consumption levels, as well as the packet delay rate across the attack connection, are then computed. The subsections that follow will detail the tests, measurements, and results.

5.1. Testbed

A laptop with the following specifications was utilized for the tests in an emulated environment: Windows 11, Core i5 11th generation, 16 GB RAM, two virtual machines running Ubuntu-16.04.2-serveramd64, and Mininet 2.2.2. Figure 4 depicts the testbed topology. The RYU controller was installed on a virtual system with 2 core CPU and four gigabytes of RAM. Mininet's topology consists of 5 switches and 8 hosts. The Mininet was housed in a virtual computer with two core and four gigabytes of RAM.

A maximum pace of 100 packets per second were allowed in each attacking host's setup, or 4.2 kilobytes per second per attacking host. Despite the Scapy library's delayed creation and transmission of spoof packets, it is sufficient in this testbed to observe the effects of a DDoS attack. The measurements were taken

over a 4-minute period, with the attack launched by a single host using a spoofed internet protocol (IP) address. It's worth noting that the DDoS launched at the very first second of the test. Table 1 summarizes the details to provide an overview of the average rate of packets arriving at the controller for processing in each attack scenario. Figure 4(a) presents the CPU usage of the RYU controller in two scenarios: before the DDoS attack and during the attack, while Figure 4(b) focuses on the network utilization of the RYU controller before the DDoS attack and during the attack. Figure 5 presents a comparison of CPU and RAM usages in the Mininet (data plane side).

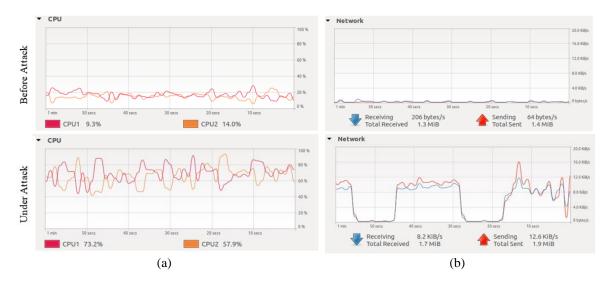


Figure 4. Before and under attack for control plane of RYU controller (a) RYU CPU usage and (b) RYU network utilize

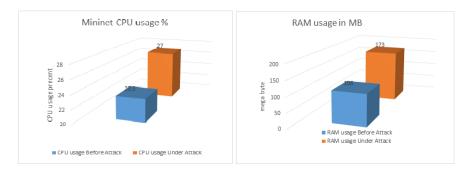


Figure 5. CPU and RAM of mininet before and under attack

| Table 1. Average link | latency between h1 an | d h2 |
|-----------------------|-----------------------|------|
| Before attack | Under attack | |
| 0.5334 ms | 4.7465 | |

5.2. Metrics

Both the CPU utilization and the data transfer rate on the control plane were measured before and after the attack to demonstrate the impact of the attack on a network. On the other side, link delay between data plane hosts was estimated, both before and during DDoS attacks, along with the RAM and CPU consumption. Mininet CPU use, link delay, and RAM usage in the data plane are the measures used to gauge attack performance.

Additionally, the control plane's network bandwidth utilization and CPU usage by the RYU controller. The Linux monitor tool was employed to measure CPU usage, and it takes a snapshot of the percentage of CPU usage at any given time every two seconds and stores it. The duration required for a packet to traverse a specific link is referred to as link latency. In this research paper, link latency was measured between hosts 1 and 2 under

two conditions: before and during the attack as shown in Figure 6. The authors devised a strategy for measuring link latency using the Mininet CLI monitor, as shown in Figure 7.

Due to the flow rule that was previously instantiated at the switch, once the packet had completed its journey across the link, it was returned to the controller. The link latency was calculated by calculating the total time required for the UDP packet to complete a round trip between hosts 1 and 2. Figure 8 explain the SDN topology using the RYU controller, the RYU controller acts as the central management entity that controls the behavior of the network switches.

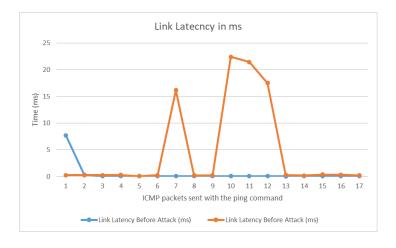


Figure 6. Link latency in ms at data plane (mininet)

| mininet> h1 ping h2 | mininet> h1 ping h2 |
|--|--|
| PING 10.0.0.2 (10.0.0.2) 56(84) bytes of data. | PING 10.0.0.2 (10.0.0.2) 56(84) bytes of data. |
| 64 bytes from 10.0.0.2: icmp_seq=1 ttl=64 time=7.69 ms | 64 bytes from 10.0.0.2: icmp seq=20 ttl=64 time=0.243 ms |
| 64 bytes from 10.0.0.2: icmp_seq=2 ttl=64 time=0.337 ms | 64 bytes from 10.0.0.2: icmp seq=21 ttl=64 time=0.259 ms |
| 64 bytes from 10.0.0.2: icmp_seq=3 ttl=64 time=0.063 ms | 64 bytes from 10.0.0.2: icmp seq=22 ttl=64 time=0.248 ms |
| 64 bytes from 10.0.0.2: icmp_seq=4 ttl=64 time=0.070 ms | 64 bytes from 10.0.0.2: icmp seq=23 ttl=64 time=0.303 ms |
| 64 bytes from 10.0.0.2: icmp_seq=5 ttl=64 time=0.066 ms | 64 bytes from 10.0.0.2: icmp seq=24 ttl=64 time=0.075 ms |
| 64 bytes from 10.0.0.2: icmp_seq=6 ttl=64 time=0.068 ms | 64 bytes from 10.0.0.2: icmp seg=25 ttl=64 time=0.264 ms |
| 64 bytes from 10.0.0.2: icmp_seq=7 ttl=64 time=0.060 ms | 64 bytes from 10.0.0.2: icmp seq=9 ttl=64 time=16159 ms |
| 64 bytes from 10.0.0.2: icmp_seq=8 ttl=64 time=0.086 ms | 64 bytes from 10.0.0.2: icmp seq=26 ttl=64 time=0.221 ms |
| 64 bytes from 10.0.0.2: icmp_seq=9 ttl=64 time=0.076 ms | 64 bytes from 10.0.0.2: icmp seq=27 ttl=64 time=0.234 ms |
| 64 bytes from 10.0.0.2: icmp_seq=10 ttl=64 time=0.063 ms | 64 bytes from 10.0.0.2: icmp seg=5 ttl=64 time=22407 ms |
| 64 bytes from 10.0.0.2: icmp_seq=11 ttl=64 time=0.085 ms | 64 bytes from 10.0.0.2: icmp seq=6 ttl=64 time=21424 ms |
| 64 bytes from 10.0.0.2: icmp_seq=12 ttl=64 time=0.062 ms | 64 bytes from 10.0.0.2: icmp seq=10 ttl=64 time=17484 ms |
| 64 bytes from 10.0.0.2: icmp_seq=13 ttl=64 time=0.052 ms | 64 bytes from 10.0.0.2: icmp seq=28 ttl=64 time=0.288 ms |
| 64 bytes from 10.0.0.2: icmp_seq=14 ttl=64 time=0.101 ms | 64 bytes from 10.0.0.2: icmp seq=29 ttl=64 time=0.196 ms |
| 64 bytes from 10.0.0.2: icmp_seq=15 ttl=64 time=0.071 ms | 64 bytes from 10.0.0.2: icmp seq=30 ttl=64 time=0.351 ms |
| 64 bytes from 10.0.0.2: icmp_seq=16 ttl=64 time=0.061 ms | 64 bytes from 10.0.0.2: icmp seq=31 ttl=64 time=0.312 ms |
| 64 bytes from 10.0.0.2: icmp_seq=17 ttl=64 time=0.056 ms | 64 bytes from 10.0.0.2: icmp_seq=32 ttl=64 time=0.223 ms |
| Refore Attack | After Attack |

Figure 7. Link latency calculation by PING command between host 1 and host 2

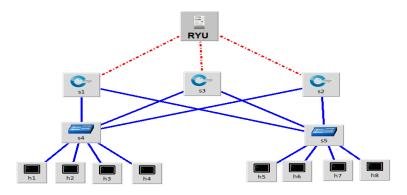


Figure 8. Evaluation of SDN topology using RYU controller

5.3. Results

Observed that the measurement of link latency was conducted by monitoring the Mininet terminal. However, we discovered that the application stopped recording data after 12 seconds of simulation time due to a controller jam, as indicated by the blue and orange vertical lines on the right side of Figure 4. This served as

an indicator of when the controller started to malfunction. Notably, the simulation with four attacking hosts resulted in the shortest time before the application failure, which was approximately 4 minutes.

Examining the CPU usage in each of the three scenarios shown on the left of Figures 4 and 5, we observed that as the number of attacking hosts increased, the fluctuation in CPU consumption decreased. It tended to reach a constant maximum value, effectively saturating the controller. The scenario with the DDoS attack caused the controller's CPU consumption to exceed 60%, representing the most severe case in terms of resource utilization. Furthermore, the time taken for a UDP packet to traverse the link between host 1 and host 2 increased significantly under the attack, as shown in Figures 6 and 7. This signifies a substantial delay in packet circulation within the network, particularly when the attack is actively deployed.

6. **CONCLUSIONS AND FUTURE WORK**

Finally, this paper concentrated on the use of a flooding attack in an SDN network to disrupt the services of a RYU. The results highlighted the severe impact of the attack on CPU consumption, network performance and RAM usage in the controller and network traffic delay. The findings showed that as the number of attacking hosts increased, so did the controller's workload. When DDoS attack in an SDN launch was determined by the RYU controller's available resources and the attack power in terms of the number of compromised hosts.

In the future, developing an intrusion detection system (IDS) within the controller as a future project. This IDS would classify network traffic and detect anomalies using software-based techniques and machine learning algorithms. Overall, this study contributes to a better understanding of the vulnerabilities and consequences of DDoS attacks in SDN architectures, emphasizing the importance of robust defense mechanisms to protect the controller while maintaining network integrity and performance.

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