Multi-level cyber security system for VANET

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Article Info	ABSTRACT			
<i>Article history:</i> Received Dec 4, 2019 Revised Feb 6, 2020 Accepted Feb 19, 2020	Recently, the cyber security of vehicular ad-hoc network (VANET) includin two practicable: car-to-car and car-to-infrastructure has been considered di to importance. It has become possible to keep pace with the development the world, where the safety of people is a priority in the development technology in general and in particular in the field of VANET. In this pape a cyber security system for VANET has been proposed to tackle the DC			
<i>Keywords:</i> C2C C2I Cyber security DoS attack VANET	attacks. The proposed system includes three security levels. Firstly, the registeration level that ask vehicles to be registered in the system, in which the network exclude the unrigestered ones. Secondly, the authentication level that checks the vehicles before joining the network. This is done by applying a proposed algorithm that considers the hash function and factory number. Thirdly, the proposed system ables to detect the DOS attack by any involved vehicle using the monitoring algorithm that allocate such vehicle to be excluded from the network. The obtained results show the efficient performance of the proposed system in managing the security of the VANET network. The multi-level system increases the security of the network, in which the attacks can be eliminated.			
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1. INTRODUCTION

Road transport focuses on people because the safety of them is the one of concerns in building a VANET network. In many countries, licensing requirements and safety regulations guarantee the separation of these two sectors. Road traffic is made by car or truck that can be targeted by DoS attackers as a prevalent in VANET. These attacks lead to damage of the safety of stirrups, and made them take a path that may lead to harm them and enter the maze of traffic congestion [1, 2].

Different studies and research work in the field of security in VANET had been presented to tackle the araised problems in terms of communication, data transmission and people safety. In [3, 4], the authors focused on Service Denial (SD) attacks that prevents end-users receiving the right data at the right moment. They analyzed SD attacks, behavior and effects on the network using various analytical models to detect an efficient answerusing different steps. Step one was building the simulator of VANET that analyzed the behavior of the network under this attack. Step two compared the proposed network with a normal network without attack to grasp the network's behavior under this type of attacks. Step three used many mathematical methods to predict if a car was under an attack or not. Authors have observed that logistic regression and neural network are better than Mean Squared Error (MSE), Root Mean Square (RMS) and Means Absolute Value (MAV) to analyze information and to predict attacks. Each parameter from the simulation was studied using the logistic regression to read the effects in details. In [5, 6], VANET turned any shared node into a wireless router network or car to cover a range of 100 to 300 meters as a base for establishing a network

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with a wide group. The authors presented an absolute overview and discussion in VANET to give a rich survy. In [7], a malicious and irrelevant packet detection algorithm were proposed, used to analyze and find the Denial of Service attack. It also lessened the overhead delay in the information analysis and processing, which increased the connect in coming times. The authors of [8, 9] suggested a reaction technique against Denial-of-Service attacks in VANET. They believed that these contributions were very useful for determining the Denial-of-Service attack reaction issue on VANETs. In [10], the authors classified the VANET system based on its capability, following attackers can be responsible for the Denial of Service (DOS) attack. In [11], authors presented a set of possible denial-of-service attacks. The first level talked about attacks that occur between Vehicles to Vehicles (V2V). It pointed out that the attacker sends false messages to the victim to be occupied and separated from the service in the network. On the other side, the intruder sends messages to the RSU running busy without contacting other Vehicles to Infrastructure (V2I).

As a result, the literary study of some researchers regarding DoS attacks leads to considerd a lightweight protocol based on Vasudev work. The proposed system differs in terms of construction, phases, and handling of DoS attacks, unlike what Vasudev did. The proposed system supports two different types of communication, Car-to-Car (C2C) and Car-to-Infrastructure (C2I).

2. PROPOSAL SYSTEM

As mentioned earlier, the proposed system considers the security at cyber level for VANET. The proposed system can be divided into different sub-sections as follows.

2.1. Proposal system structure

Figure 1, Shows the the proposed schema that uses the lightweight protocol due to limitation of storage memory employing the technology of time difference. This schema explains the three security levels of the proposed system: registeration, authentication and DOS attack detector for four cars as a prototype. At the first level, each car (A, B, C, D) request individual keys from the server as a registeration process. The server responds requests by generated keys. The lightweight authentication protocol is implemented in the second level by exhchanging keys with nighbour vehicles. If the system detects unauthenticated vehicle, it is excluded from the network. The third level explains the detection of DOS attack by one of included vehicle. The detected attacked vehicle is excluded from the network.



Figure 1. Proposal system structure

2.2. Levels of the proposed system

As mentioned earlier, three levels are considered in the proposed cyber security system. These levels can be explaining in the next sub-sections.

2.2.1. Level One: registration

In the registration phase, each vehicle sends the key request to the server to be registered within the system through the manufacturers or traffic. At the other hand, the server stores all information, sent from the vehicles, within the TAMPER-PROOF DEVICE (TPD) [12-15]. After that, each vehicle works to store its key also within its own TPD. This is done through a safe path that is the result of applying the proposed protocol works to secure it as hown in Figures 2. The working steps of the proposed registeration algorithm is shown as:

a) Every car: (car_a ,car_b ,car_c ,car_d) selects and Identification (ID): (ID_car_a , ID_car_b ,ID_car_c ,ID_car_d) and PassWords (PW): (PW_car_a ,PW_car_b ,PW_car_c ,PW_car_d) and they generate values: (R_car_a ,R_car_b ,R_car_c ,R_car_d). For ensuring security, all cars compute the following:

Car_a:	Car_c:
$A_a = h [ID_car_a PW_car_a] \bigoplus R_car_a$	$A_c = h [ID_car_c PW_car_c] \bigoplus R_car_c$
$B_a = h [ID_car_a A_a] \bigoplus R_car_a$	$B_c = h [ID_car_c A_c] \bigoplus R_car_c$
$C_a = h [PW_car_a B_a] \bigoplus R_car_a$	$C_c = h [PW_car_c B_c] \bigoplus R_car_c$
Req.id = $B_a \bigoplus C_a$	Req.id = $B_c \bigoplus C_c$
Car_b:	Car_d:
$A_b = h [ID_car_b PW_car_b] \bigoplus R_car_b$	$A_d = h [ID_car_d \parallel PW_car_d] \bigoplus R_car_d$
$B_b = h [ID_car_b A_b] \bigoplus R_car_b$	$B_d = h [ID_car_d \parallel A_d] \bigoplus R_car_d$
$C_b = h [PW_car_b B_b] \bigoplus R_car_b$	$C_d = h [PW_car_d \parallel B_d] \bigoplus R_car_d$
Req.id = $B_b \bigoplus C_b$	Req.id = $B_d \bigoplus C_d$

Then encryption of Req.id, which is the requested key, is sent to Car_Server through insecure path. b) After receiving the Car-Server the requests, it selects the ID, PW and time response (T_S): (ID_S, PW_S, T_S) and generate values (R_S). For ensuring security, the Car_Server computes:

Ns = h [$ID_car_S \parallel T_S] \oplus R_car_S$

Key_car_i = h [ID_car_S \parallel Req.i \parallel Ns] \bigoplus R_car_S

Then, the TPD_Server Stores the result as [Key_car_i]. The encrypted Key_car_i is sent to all cars (Car_a, Car_b, Car_c, Car_d) through insecure path.

c) After receiving the keys of cars (Car_a , Car_b , Car_c , Car_d), the decryption process is done for Key_car_i. Then TPD_Car Stores its own key (Key_car_a , Key_car_b , Key_car_c , Key_car_d).

Figure 2, Explain the algorithm of registeration process for car A as example. The other cars follows the same steps.

Car_a	Car_Server
Choose ID_car_a, PW_car_a.	
Random R_car_a .	
Computes	
$A_a = h [ID_car_a \parallel PW_car_a] \oplus R_car_a$	L
$B_a = h [ID_car_a \parallel A_a] \bigoplus R_car_a$	
$C_a = h [PW_car_a \parallel B_a] \bigoplus R_car_a$	
$\operatorname{Req.i} = \operatorname{B}_{a} \bigoplus \operatorname{C}_{a}$	
ENC [Req.i]	DEC [Pagi]
	Chaosa ID car S DW car S T S
	Choose ID_caf_5, Pw_caf_5, 1_5.
	Kandom R_car_S.
	Computes
	$N_s = h [ID_car_S \parallel T_S] \bigoplus R_car_S$
	$Key_car_i = h [ID_car_S Req.i N_s] \bigoplus R_car_S$
	TPD_Server Storage (Key_car_i)
	ENC [Key car i]
DEC [Key_car_i]	
TPD_car_a Storage	
[Key_car_a]	

Figure 2. Vehicle registration A

2.2.2. Level two: authentication

During the authentication phase, any vehicle inside the network can communicate with another vehicle in specific range using the technique of Dedicated Short Range Communication (DSRC) [16-20]. The authentication process gives the licence for cars to communicate. For example, when the vehicle A is authenticated with vehicle B, vehicle A can send its key to vehicle B. While, vehicle B can send its key to vehicle A and thus both cars now have the other vehicle key (Car_a = {Key_car_b} and Car_b = {Key_car_a}). Vehicle A sends the key of vehicle B to the Server_Car and vehicle B is doing the same for A. Now, the presence of the A and B vehicles is checked if the key for each vehicle registered within the TPD server. If the result is positive, then such car is treated as a licensed vehicle. At this stage, it cannot be confirmed that the vehicle is harmful or not, only, the key is checked in the server after which the vehicle is communicated. Figure 3, Explains the procedure of authentication between car A and B as an algorithm.



Figure 3. Authentication between Car A & B

2.2.3. Level three: communication and DOS attack detection

In this level, the proposed protocol works to secure data, discover the attack [21-26]. Know the harmful vehicle after it has passed through the registration and authentication stages. The communication and data transfer process is performed between vehicle A and as shown in Figure 4, That viewed the diagrammatic of the Car-to-Car.

a) Car_a, selects: (Welcome, Key_car_a, Sending time T_a) and generates value of (U_car_a), then for ensuring security, the car computes:

Car_a : $X_a = h [Key_car_a || T_a] \bigoplus U_car_a$ $Y_a = h [X_a || Key_car_a] \bigoplus U_car_a$ $Z_a = IN \bigoplus Key_car_a \bigoplus X_a \bigoplus Y_a \bigoplus T_a$

The encryption of Req and is sent to Car_b through insecure path in the following form:

 T_a , Req = ENC H_a (Z_a), where ENC is stand for encryption.

b) After receiving the Car_b the request, it performs the decryption for Req (Req = DEC H_b (Z_a)), where DEC is stand for decryption. The time difference verification is computed depending on ($\Delta T_b \leq T_b - T_b$

 T_a). The car selects: (Welcome, Key_car_b, Sending time T_b) and generates value of (U_car_b), then for ensuring security, the Car computes:

Car_b :

 $\begin{array}{l} X_{b} = h \; [\; Key_car_b \; \| \; T_b \;] \bigoplus U_car_b \\ Y_{b} = h \; [\; X_{b} \; \| \; Key_car_b \;] \bigoplus U_car_b \\ Z_{b} = IN \bigoplus Key_car_b \bigoplus X_{b} \bigoplus Y_{b} \bigoplus \Delta T_{b} \end{array}$

The encryption of Rep and sent to Car_a through insecure path in the following form:

 $Rep = ENC H_b (Z_b)$, Sending time : T_b'

c) In the Car_a side and after receiving the request, the decryption is done for Rep (Rep = DEC H_a (Z_b)). The time difference verification is evaluated depending on ($\Delta T_a \leq T_a' - T_b'$), after that the communication process is performed.



Figure 4. Communication and attack detection

2.3. GUI of the proposed system

This section discusses the graphical user interface shown in Figure 5, Which represents the practical aspect in this paper. The language C# is used in the construction of the network as well as the security aspect, which represents the security of data transmission between vehicles. The server is built from the following parts: First part is the environment that contains a group of the vehicles (A, B, C, D) where A and B are normal vehicles while D represents the attack of DoS and C has malfunction of its system during the authentication phase. The second part represents the network infrastructure and is the presence of the server that contains the information about vehicles of the network. The third part is the tools through which the three security levels are implemented: registration, authentication, as well as the communication and attack detection. The last part is a window of events in which the results of the communication phases appear in the proposed protocol.





Figure 5. GUI of the proposed sytsem

3. **RESULTS**

It is well known now that the considered prototype includes four cars. In order to test the performance of the proposed security system, three case studies has been considered.

3.1. Case study one

In the registration level of vehicles (A, B, C, D), they send their information, which represents a request for a key for each vehicle in the network to the server. After accepting the registration request of vehicles, the server works to send a key to each vehicle that its information is registered by the server (See Figures of 6-9). Note that there is no vehicle can access the network before registration.

Registration Vehicle [A]	А		Vehicle A → 70U7zgllbXA=	Registration	В	Vehicle B → nC6IKMQlgrs=
	•	s	Server A → 70U7zgllbXA=	Vehicle [B]	•	

Figure 6. Vehicle registration A

Figure 7. Vehicle registration B



Figure 8. Vehicle registration C

Figure 9. Vehicle registration D

3.2. Case study two

After the registration phase, the authentication phase is performed where the vehicles are authenticated with each other and the server according to the mechanism of the proposed authentication algorithm. Authentication process between vehicles A and B as well as D and B is normal and successed without problems as shown in Figures 10 and 11. With regard to vehicle authentication of vehicles C and B are unsuccessful due to a malfunction of the vehicle C or a change in vehicle information after the registration process. Therefore, the communication between car C and B has not been authenticated, and the vehicle C is blocked and banned from the grid as shown in Figure 12.

Authentication Vehicle A&B	A Key B		Key A B	
Check	Vehicle Server			
Key [A & B]	Server \rightarrow Existing			
Authentication Resulting		Successful		

Authentication	D	\rightarrow	Key D		
Vehicle D&B	Key B	•	В		
Check	Vehicle Server				
Key [D & B]	Server \rightarrow Existing				
Authentication Resulting		Successful			

Figure 10. Authenticaton between vehicle A & B

Figure 11. Authenticaton between vehicle D & B

Authentication	С		Key C		
Vehicle C&B	Key B	•	В		
Check	Vehicle Server				
Key [C & B]	Server \rightarrow Not Existing				
Authentication Resulting		Unsuccessful			

Figure 12. Authenticaton between Vehicle C & B

3.3. Case study three

Here we address two cases. The first case is a normal transmission process where a welcome message (hello) is sent from the vehicle A after encrypted by hash function to vehicle B. Thus, the vehicle works to calculate the time difference of the number of messages received and make sure their integrity and decoding. The message is sent back to the vehicle A also calculates the time difference of the received messages to verify the integrity of the message to be a successful picture transmission without malfunction or attack during the transmission process. Thus it indicates the strength and durability of the protocol in force as shown in Figure 13.

The second case is the process of detecting an attack by sending an unspecified large set of messages built by the DoS attack in vehicle D. It represents the role of the DoS attack by sending a number of messages that indicate (accident at location Y) to the victim vehicle B and making it takes another road that might be congested traffic. This attack is discovered by calculating the time difference of the number of messages, sent from vehicle D inside vehicle B, and thus identify the attacking vehicle. The proposed system blockes vehicle D information inside the server for disconnecting it from the network. The proposed system does not allow the attacker to enter the network again as shown in Figure 14.

		Hello , T.a	Check
C O M U N	Vehicle A	D1+ksRULieXGJRbOVeouO BNqF9YY3Yd0q5u8CVW/E 4vGJRbOVeouOF9qHVagn Q6E20Xlhjdh3T5Yt+P+N+N JHHSI+rULIVw	$\Delta T_b \leq T.b - T.a$
Ι	Check	Hello , T.b'	
C A T I O N	$\Delta T_a \le T.a' - T.b'$	↓ D2+ksRULieXGJRbOVeouO BNqF9YY3Yd0q5u8CVW/E 4vGJRbOVeouOF9qHVagn Q6E20Xlhjdh3T5Yt+P+N+N JHHSI+rULIVw	Vehicle B
	True	Successful	True

Figure 13. Communication from vehicle A to vehicle B (NORMAL)

		Accident at location Y, T.d, T.d', T.d"	Check
C O M U	Vehicle D (attack)	EpqpDSIHop4MFYOL7NHPR B168Hj6aWiTHTDn/LadffFST Y5vK6dJEGNM7zwb8WPNW ELSTQ9JZWDN1877HIEEcslo y72SeYVLJBxVtZco0fAICMIR KV3L3HsI+rULIVW	ΔT _b ≤ T.b -T.d
N	Check	Accident at location Y, T.b'	
I C A T I O N	∆T _d ≤T.d' -T.d'	EpqpDSIHop4MFYOL7NHPR B168Hj6aWITHTDn/LadffFST Y5vK6dJEGNM7zwb8WPNW ELSTQ9JZWDN1877HIEEcslo y72SeYVLJBxVtZco0fAICMIR KV3L3HsI+rULIVW	Vehicle B
	True	Unsuccessful	false

Figure 14. Communication from vehicle D (DoS attack) to vehicle B (ATTACK)

4. CONCLUSION

In this paper, we have proposed a lightweight protocol for multi-level cyber security in VANET. The proposed protocol consisted of three levels, each of which works to maintain network security, from attacks that are related to DoS attacks to reach the required safety. These levels were registeration, authentication as well as communication and attach detection. The proposed levels worked as obstacles to prevent the DoS attaches. Even if the the attached vehicle passes the registeration and authentication levels, the third level can detect it from its behaivour inside the VANET. The obtained results showed the efficiency in performance of the peoposed system in detecting the attckes. This was concluded by considering different case studies.

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